

The maximum segment product

Given integer array A of length N we are asked to compute the maximum segment product of A . More formally, for natural numbers p and q and integer r our program must establish

$$R : \quad r = \langle \uparrow p, q : 0 \leq p \leq q \leq N : M.p.q \rangle$$

where, for natural i , M is defined as

$$M.p.q = \langle \Pi i : p \leq i < q : A.i \rangle$$

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We investigate our standard invariant, P , which is the conjunction of

$$Pr : \quad r = \langle \uparrow p, q : 0 \leq p \leq q \leq n : M.p.q \rangle$$

$$Pn : \quad 0 \leq n \leq N$$

which we initially establish with $n, r := 0, 1$.

We progress to termination by incrementing n by one. The proofs of termination and maintenance of Pn under $n := n + 1$ are standard and trivial and we thus omit them.

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We observe

$$\begin{aligned} & \llbracket \text{Context: } P \wedge n \neq N \\ & \quad \langle \uparrow p, q : 0 \leq p \leq q \leq n + 1 : M.p.q \rangle \\ & = \quad \{ \text{split off } q = n + 1, \text{ requires } n + 1 \leq N ; Pr \} \\ & \quad r \uparrow \langle \uparrow p : 0 \leq p \leq n + 1 : M.p.(n + 1) \rangle \\ & = \quad \{ \text{assuming } (n := n + 1).Ps \} \\ & \quad r \uparrow s \\ & \rrbracket \end{aligned}$$

We strengthen P with

$$Ps : \quad s = \langle \uparrow p : 0 \leq p \leq n : M.p.n \rangle$$

which we establish initially with $s := 1$.

At this point we have

$$\begin{aligned}
& | [\textit{Context}: P \wedge n \neq N \\
& \quad \{Pr \wedge (n := n+1).Ps\} \quad r := r \uparrow s \quad \{(n := n+1).Pr\} \\
&]]
\end{aligned}$$

* * *

We turn our attention to maintenance of Ps . We observe

$$\begin{aligned}
& | [\textit{Context}: P \wedge n \neq N \\
& \quad \langle \uparrow p: 0 \leq p \leq n+1: M.p.(n+1) \rangle \\
& = \quad \{ \text{split off } p = n+1, \text{ requires } n+1 \leq N \} \\
& \quad \langle \uparrow p: 0 \leq p \leq n: M.p.(n+1) \rangle \uparrow M.(n+1).(n+1) \\
& = \quad \{ \text{def of } M \} \\
& \quad \langle \uparrow p: 0 \leq p \leq n: M.p.n * A.n \rangle \uparrow 1 \\
& = \quad \{ \text{case analysis} \} \\
& \quad \left\{ \begin{array}{l} \langle \langle \uparrow p: 0 \leq p \leq n: M.p.n \rangle * A.n \rangle \uparrow 1, \text{ if } 0 \leq A.n \\ \langle \langle \downarrow p: 0 \leq p \leq n: M.p.n \rangle * A.n \rangle \uparrow 1, \text{ if } A.n \leq 0 \end{array} \right. \\
& = \quad \{ Ps; Pt \} \\
& \quad \left\{ \begin{array}{l} (s * A.n) \uparrow 1, \text{ if } 0 \leq A.n \\ (t * A.n) \uparrow 1, \text{ if } A.n \leq 0 \end{array} \right. \\
&]]
\end{aligned}$$

Once again we have strengthened P , this time with

$$Pt: \quad t = \langle \downarrow p: 0 \leq p \leq n: M.p.n \rangle$$

We now have

$$\begin{aligned}
& | [\textit{Context}: P \wedge n \neq N \\
& \quad \{Ps \wedge Pt\} \\
& \quad \mathbf{if} \ 0 \leq A.n \ \rightarrow \ s \ := \ (s * A.n) \uparrow 1 \\
& \quad \quad \square \ A.n \leq 0 \ \rightarrow \ s \ := \ (t * A.n) \uparrow 1 \\
& \quad \mathbf{fi} \\
& \quad \{(n := n+1).Ps\} \\
&]]
\end{aligned}$$

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Let us now address the maintenance of Pt . We observe

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| [ Context: P ∧ n ≠ N
  ⟨↓ p: 0 ≤ p ≤ n + 1: M.p.(n + 1)⟩
=   { split off p = n + 1 }
  ⟨↓ p: 0 ≤ p ≤ n: M.p.(n + 1)⟩ ↓ M.(n + 1).(n + 1)
=   { def of M }
  ⟨↓ p: 0 ≤ p ≤ n: M.p.n * A.n⟩ ↓ 1
=   { case analysis }
  { (⟨↓ p: 0 ≤ p ≤ n: M.p.n⟩ * A.n) ↓ 1, if 0 ≤ A.n
    (⟨↑ p: 0 ≤ p ≤ n: M.p.n⟩ * A.n) ↓ 1, if A.n ≤ 0 }
=   { Ps; Pt }
  { (t * A.n) ↓ 1, if 0 ≤ A.n
    (s * A.n) ↓ 1, if A.n ≤ 0 }
] ]

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This gives us

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| [ Context: P ∧ n ≠ N
  {Ps ∧ Pt}
  if 0 ≤ A.n → t := (t * A.n) ↓ 1
  [] A.n ≤ 0 → t := (s * A.n) ↓ 1
  fi
  {(n := n + 1).Pt}
] ]

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Our solution is thus

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n, r, s, t := 0, 1, 1, 1
; do n ≠ N →
  if 0 ≤ A.n → s, t := (s * A.n) ↑ 1, (t * A.n) ↓ 1
  [] A.n ≤ 0 → s, t := (t * A.n) ↑ 1, (s * A.n) ↓ 1
  fi
; r, n := r ↑ s, n + 1
od

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