

## Multiplication

Given  $G$  defined by

$$G.x.0 = 0$$

$$G.x.y = G.(x * 2).(y \text{ div } 2) \quad \text{if } y \bmod 2 = 0$$

$$G.x.y = x + G.x.(y - 1) \quad \text{if } y \bmod 2 = 1$$

we are asked to derive, for integers  $A$  and  $B$ , a program to compute  $G.A.B$ .

\* \* \*

As invariant we choose  $P$  given by

$$Py : \quad 0 \leq y$$

$$Pr : \quad G.A.B = r + G.x.y$$

which we initially establish with  $x, y, r := A, B, 0$ .

The postcondition is

$$R : \quad r = G.A.B$$

\* \* \*

The repetition can terminate whenever  $LHS.R = RHS.Pr$  :

$$\begin{aligned} r &= r + G.x.y \\ &= \{ \text{arithmetic} \} \\ 0 &= G.x.y \\ \Leftarrow & \{ \text{properties of } G \} \\ y &= 0 \end{aligned}$$

and therefore  $y \neq 0$  is an acceptable guard. Hence we are heading for a program of the form

```
 $x, y, r := A, B, 0$ 
; do  $y \neq 0 \rightarrow$ 
    “Decrease  $y$  under invariance of  $P$  ”
od
```

\* \* \*

We observe

$$\begin{aligned} & \llbracket \textit{Context: } P \wedge y \neq 0 \wedge y \bmod 2 = 0 \\ & \quad r + G.x.y \\ & = \quad \{ y \bmod 2 = 0, \text{ property of } G \} \\ & \quad r + G.(x * 2).(y \mathbf{div} 2) \\ & \rrbracket \end{aligned}$$

and

$$\begin{aligned} & \llbracket \textit{Context: } P \wedge y \neq 0 \wedge y \bmod 2 = 1 \\ & \quad r + G.x.y \\ & = \quad \{ y \bmod 2 = 1, \text{ property of } G \} \\ & \quad r + x + G.x.(y - 1) \\ & \rrbracket \end{aligned}$$

and the resulting program is

```
x, y, r := A, B, 0
; do y ≠ 0 →
  if y mod 2 = 0 → x, y := x * 2, y div 2
  [] y mod 2 = 1 → r, y := r + x, y - 1
  fi
od
```

The proofs of termination and maintenance of  $Py$  under  $y := y - 1$  and  $y := y \mathbf{div} 2$  are standard and we thus omit them.

*E. Emmanuel Macaulay*  
eric@mathmeth.com

22 January 2007