

fusc

In *Programming: The Derivation of Algorithms* Anne Kaldewaij gives us the recurrence

$$\text{fusc.0} = 0, \text{fusc.1} = 1$$

$$\text{fusc.(2 * n)} = \text{fusc.n}$$

$$\text{fusc.(2 * n + 1)} = \text{fusc.n} + \text{fusc.(n + 1)} \quad \text{for } n \geq 0$$

and asks us to derive a program for the computation of fusc.N , $N \geq 0$.

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The postcondition of our program is

$$R: \quad r = \text{fusc.N}$$

Our next task is to choose an invariant. The shape of the recurrence suggests as invariant the conjunction of

$$\langle \text{fusc.n} + \text{fusc.(n + 1)} \rangle$$

$$Pn: \quad 0 \leq n \leq N$$

where

$$\langle x \rangle \equiv \text{fusc.N} = x$$

Kaldewaij hints that we should compute fusc.78 , so let us do so. We can compute this in either a depth-first or breadth-first manner. We choose the latter as it is simpler (we do not have to store fusc values for later use).

$$\begin{aligned} & \text{fusc.78} \\ = & \{ \} \\ & \text{fusc.39} \\ = & \{ \} \\ & \text{fusc.19} \quad + \quad \text{fusc.20} \\ = & \{ \} \\ & \text{fusc.9} \quad + \quad \text{fusc.10} \quad + \quad \text{fusc.20} \\ = & \{ \} \end{aligned}$$

$$\begin{aligned}
& fusc.9 + 2 * fusc.10 \\
= & \{ \} \\
& fusc.4 + fusc.5 + 2 * fusc.10 \\
= & \{ \} \\
& fusc.4 + 3 * fusc.5 \\
= & \{ \} \\
& fusc.2 + 3 * fusc.5 \\
= & \{ \} \\
& 4 * fusc.2 + 3 * fusc.3 \\
= & \{ \} \\
& 4 * fusc.1 + 3 * fusc.3 \\
= & \{ \} \\
& 7 * fusc.1 + 3 * fusc.2 \\
= & \{ \} \\
& 10
\end{aligned}$$

This calculation suggests we generalise the invariant to

$$P : \langle a * fusc.n + b * fusc.(n + 1) \rangle$$

which we initially establish with $n, a, b := N, 1, 0$. We are heading for a program of the form

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n, a, b := N, 1, 0
; do n ≠ 0 →
    “Decrease n under invariance of P ”
od

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The shape of $fusc$ suggests we investigate two cases: $even.n$ and $odd.n$. We observe

$$\begin{aligned}
& [| \textit{Context: } even.n \\
& \quad a * fusc.n \quad + \quad b * fusc.(n + 1) \\
& = \quad \{ \textit{property of } fusc ; \textit{algebra} \} \\
& \quad a * fusc.(n \mathbf{div} 2) \quad + \quad b * fusc.((n + 1) \mathbf{div} 2) \quad + \quad b * fusc.(((n + 1) \mathbf{div} 2) + 1) \\
& = \quad \{ (n + 1) \mathbf{div} 2 = n \mathbf{div} 2 ; \textit{follows from } odd.(n + 1) \} \\
& \quad a * fusc.(n \mathbf{div} 2) \quad + \quad b * fusc.(n \mathbf{div} 2) \quad + \quad b * fusc.((n \mathbf{div} 2) + 1) \\
& = \quad \{ \textit{algebra} \} \\
& \quad (a + b) * fusc.(n \mathbf{div} 2) \quad + \quad b * fusc.((n \mathbf{div} 2) + 1) \\
& = \quad \nabla \ a, n \ := \ a + b, n \mathbf{div} 2 \ \nabla \\
& \quad a * fusc.n \quad + \quad b * fusc.(n + 1) \\
& |]
\end{aligned}$$

and

$$\begin{aligned}
& [| \textit{Context: } odd.n \\
& \quad a * fusc.n \quad + \quad b * fusc.(n + 1) \\
& = \quad \{ \textit{property of } fusc ; \textit{algebra} \} \\
& \quad a * fusc.(n \mathbf{div} 2) \quad + \quad a * fusc.((n \mathbf{div} 2) + 1) \quad + \quad b * fusc.((n + 1) \mathbf{div} 2) \\
& = \quad \{ (n + 1) \mathbf{div} 2 = (n \mathbf{div} 2) + 1 ; \textit{follows from } even.(n + 1) \} \\
& \quad a * fusc.(n \mathbf{div} 2) \quad + \quad a * fusc.((n \mathbf{div} 2) + 1) \quad + \quad b * fusc.((n \mathbf{div} 2) + 1) \\
& = \quad \{ \textit{algebra} \} \\
& \quad a * fusc.(n \mathbf{div} 2) \quad + \quad (a + b) * fusc.((n \mathbf{div} 2) + 1) \\
& = \quad \nabla \ b, n \ := \ a + b, n \mathbf{div} 2 \ \nabla \\
& \quad a * fusc.n \quad + \quad b * fusc.(n + 1) \\
& |]
\end{aligned}$$

Hence our solution

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 $n, a, b := N, 1, 0$ 
; do  $n \neq 0 \rightarrow$ 
    if  $even.n \rightarrow a, n := a + b, n \mathbf{div} 2$ 
    []  $odd.n \rightarrow b, n := a + b, n \mathbf{div} 2$ 
    fi
od
;  $r := b$ 
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The proofs of termination and maintenance of Pn under $n := n \mathbf{div} 2$ are standard hence we omit them.

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