

A note on equation $x: [t \vee s; x \equiv x]$

We consider the "homogeneous" linear equation

$$(0) \quad x: [s; x \equiv x],$$

which - by definition - has ∞s as its weakest solution.

We also consider the "inhomogeneous" linear equation

$$(1) \quad x: [t \vee s; x \equiv x].$$

Now we have the following theorem.

Theorem The weakest solution of (1) is $r \vee \infty s$, where r is any solution of (1).

Proof

$r \vee \infty s$ solves (1):

$$\begin{aligned} & t \vee s; (r \vee \infty s) \\ = & \{r.c.\} \\ & t \vee s; r \vee s; \infty s \\ = & \{r \text{ solves (1)}\} \\ & r \vee s; \infty s \\ = & \{\infty s \text{ solves (0)}\} \\ & r \vee \infty s \end{aligned}$$

$r \vee \infty s$ is the weak extreme of (1):

We show that for any x

$$[x \Rightarrow r \vee \infty s] \Leftarrow [t \vee s; x \equiv x]:$$

$$\begin{aligned}
& [x \Rightarrow r \vee \infty s] \\
= & \{p.c.\} \\
& [x \wedge \neg r \Rightarrow \infty s] \\
\Leftarrow & \{ \infty s \text{ is the weak extreme of } (0) \} \{KT\} \\
& [s; (x \wedge \neg r) \Leftarrow x \wedge \neg r] \\
= & \{p.c.\} \\
& [r \vee s; (x \wedge \neg r) \Leftarrow x] \\
= & \{ r \text{ solves } (1) \} \\
& [t \vee s; r \vee s; (x \wedge \neg r) \Leftarrow x] \\
= & \{r.c. : \text{this is a macro}\} \\
& [t \vee s; r \vee s; x \Leftarrow x] \\
\Leftarrow & \{p.c.\} \\
& [t \vee s.x \equiv x] .
\end{aligned}$$

(End of Proof.)

The analogy of the above with ordinary differential equations is striking - hence the ad-hoc jargon - . I did not know the theorem - hence this note - . The above result has been triggered directly by Rutger-5, which uses a special instance of the theorem (Rutger Dijkstra may have been aware of the theorem.)

The question that remains is whether we should explore, in more depth, equations à la the above and their properties. I think we should, if only for fun.

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(Written with MB Meisterstück.)